

**NRI INSTITUTE OF INFORMATION
SCIENCE
& TECHNOLOGY BHOPAL**



**DEPARTMENT OF INFORMATION
TECHNOLOGY**

LAB MANUAL

OS LAB

(IT-501)

INFORMATION TECHNOLOGY (IT)

Department of Information Technology
Operating System (IT501) Lab Manual

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EXPERIMENT NO.1

CPU SCHEDULING ALGORITHMS

A). FIRST COME FIRST SERVE:

AIM: To write a c program to simulate the CPU scheduling algorithm First Come First Serve (FCFS)

DESCRIPTION:

To calculate the average waiting time using the FCFS algorithm first the waiting time of the first process is kept zero and the waiting time of the second process is the burst time of the first process and the waiting time of the third process is the sum of the burst times of the first and the second process and so on. After calculating all the waiting times the average waiting time is calculated as the average of all the waiting times. FCFS mainly says first come first serve the algorithm which came first will be served first.

ALGORITHM:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue

Step 3: For each process in the ready Q, assign the process name and the burst time Step

4: Set the waiting of the first process as 0 and its burst time as its turnaround time Step

5: for each process in the Ready Q calculate

a). $\text{Waiting time}(n) = \text{waiting time}(n-1) + \text{Burst time}(n-1)$ b).

$\text{Turnaround time}(n) = \text{waiting time}(n) + \text{Burst time}(n)$

Step 6: Calculate

a) $\text{Average waiting time} = \text{Total waiting Time} / \text{Number of process}$

b) $\text{Average Turnaround time} = \text{Total Turnaround Time} / \text{Number of process}$

Step 7: Stop the process

SOURCE CODE:

```
#include<stdio.h>
#include<conio.h>
main()
{
int bt[20], wt[20], tat[20], i, n;
float wtavg, tatavg;
clrscr();
printf("\nEnter the number of processes -- ");
scanf("%d", &n);
for(i=0;i<n;i++)
{
printf("\nEnter Burst Time for Process %d -- ", i);
scanf("%d", &bt[i]);
}
wt[0] = wtavg = 0;
tat[0] = tatavg = bt[0];
for(i=1;i<n;i++)
{
wt[i] = wt[i-1] +bt[i-1];
tat[i] = tat[i-1] +bt[i];
wtavg = wtavg + wt[i];
tatavg = tatavg + tat[i];
}
printf("\t PROCESS \tBURST TIME \t WAITING TIME\t TURNAROUND TIME\n");
for(i=0;i<n;i++)
printf("\n\t P%d \t\t %d \t\t %d \t\t %d", i, bt[i], wt[i], tat[i]);
printf("\nAverage Waiting Time -- %f", wtavg/n);
printf("\nAverage Turnaround Time -- %f", tatavg/n);
getch();
}
```

INPUT

Enter the number of processes -- 3
Enter Burst Time for Process 0 -- 24
Enter Burst Time for Process 1 -- 3
Enter Burst Time for Process 2 -- 3

OUTPUT

PROCESS	BURST TIME	WAITING TIME	TURNAROUND TIME
P0	24	0	24
P1	3	24	27
P2	3	27	30
Average Waiting Time--	17.000000		
Average Turnaround Time --		27.000000	

B). SHORTEST JOB FIRST:

AIM: To write a program to stimulate the CPU scheduling algorithm Shortest job first (Non- Preemption)

DESCRIPTION:

To calculate the average waiting time in the shortest job first algorithm the sorting of the process based on their burst time in ascending order then calculate the waiting time of each process as the sum of the bursting times of all the process previous or before to that process.

ALGORITHM:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: Start the Ready Q according the shortest Burst time by sorting according to lowest to highest burst time.

Step 5: Set the waiting time of the first process as $_0$ and its turnaround time as its burst time.

Step 6: Sort the processes names based on their Burt time

Step 7: For each process in the ready queue, calculate

a) $\text{Waiting time}(n) = \text{waiting time}(n-1) + \text{Burst time}(n-1)$

b) $\text{Turnaround time}(n) = \text{waiting time}(n) + \text{Burst time}(n)$

Step 8: Calculate

c) $\text{Average waiting time} = \text{Total waiting Time} / \text{Number of process}$

d) $\text{Average Turnaround time} = \text{Total Turnaround Time} / \text{Number of process}$

Step 9: Stop the process

SOURCE CODE :

```
#include<stdio.h>
#include<conio.h>
main()
{
int p[20], bt[20], wt[20], tat[20], i, k, n, temp; float wtavg,
tatavg;
clrscr();
printf("\nEnter the number of processes -- ");
scanf("%d", &n);
for(i=0;i<n;i++)
{
p[i]=i;
printf("Enter Burst Time for Process %d -- ", i);
scanf("%d", &bt[i]);

}
for(i=0;i<n;i++)
for(k=i+1;k<n;k++)
if(bt[i]>bt[k])
{
temp=bt[i];
bt[i]=bt[k];
bt[k]=temp;

temp=p[i];
p[i]=p[k];
p[k]=temp;
}
wt[0] = wtavg = 0;
tat[0] = tatavg = bt[0]; for(i=1;i<n;i++)
{
wt[i] = wt[i-1] +bt[i-1];
tat[i] = tat[i-1] +bt[i];
wtavg = wtavg + wt[i];
tatavg = tatavg + tat[i];
}
printf("\n\t PROCESS \tBURST TIME \t WAITING TIME\t TURNAROUND TIME\n");
for(i=0;i<n;i++)
printf("\n\t P%d \t\t %d \t\t %d \t\t %d", p[i], bt[i], wt[i], tat[i]);
printf("\nAverage Waiting Time -- %f", wtavg/n);
printf("\nAverage Turnaround Time -- %f", tatavg/n); getch();}
```

INPUT

Enter the number of processes -- 4
Enter Burst Time for Process 0 -- 6
Enter Burst Time for Process 1 -- 8
Enter Burst Time for Process 2 -- 7
Enter Burst Time for Process 3 -- 3

OUTPUT

PROCESS	BURST TIME	WAITING TIME	TURNARO UND TIME
P3	3	0	3
P0	6	3	9
P2	7	9	16
P1	8	16	24
Average Waiting Time --		7.000000	
Average Turnaround Time --		13.000000	

C). ROUND ROBIN:

AIM: To simulate the CPU scheduling algorithm round-robin.

DESCRIPTION:

To aim is to calculate the average waiting time. There will be a time slice, each process should be executed within that time-slice and if not it will go to the waiting state so first check whether the burst time is less than the time-slice. If it is less than it assign the waiting time to the sum of the total times. If it is greater than the burst-time then subtract the time slot from the actual burst time and increment it by time-slot and the loop continues until all the processes are completed.

ALGORITHM:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue and time quantum (or) time slice

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: Calculate the no. of time slices for each process where No. of time slice for process (n) = burst time process (n)/time slice

Step 5: If the burst time is less than the time slice then the no. of time slices =1.

Step 6: Consider the ready queue is a circular Q, calculate

a) Waiting time for process (n) = waiting time of process(n-1)+ burst time of process(n-1) + the time difference in getting the CPU from process(n-1)

b) Turnaround time for process(n) = waiting time of process(n) + burst time of process(n)+ the time difference in getting CPU from process(n).

Step 7: Calculate

c) Average waiting time = Total waiting Time / Number of process

d) Average Turnaround time = Total Turnaround Time / Number of process Step

8: Stop the process

SOURCE CODE

```
#include<stdio.h>
main()
{
int i,j,n,bu[10],wa[10],tat[10],t,ct[10],max;
float awt=0,att=0,temp=0;
clrscr();
printf("Enter the no of processes -- ");
scanf("%d",&n);
for(i=0;i<n;i++)
{
printf("\nEnter Burst Time for process %d -- ", i+1);
scanf("%d",&bu[i]);
ct[i]=bu[i];
}
printf("\nEnter the size of time slice -- ");
scanf("%d",&t);
max=bu[0];
for(i=1;i<n;i++)
if(max<bu[i])
max=bu[i];
for(j=0;j<(max/t)+1;j++)
for(i=0;i<n;i++)
if(bu[i]!=0)
if(bu[i]<=t) {
tat[i]=temp+bu[i];
temp=temp+bu[i];
bu[i]=0;
}
else {
bu[i]=bu[i]-t;
temp=temp+t;
}
for(i=0;i<n;i++){
wa[i]=tat[i]-
ct[i]; att+=tat[i];
awt+=wa[i];}
printf("\nThe Average Turnaround time is -- %f",att/n);
printf("\nThe Average Waiting time is -- %f ",awt/n);
printf("\n\tPROCESS\t BURST TIME \t WAITING TIME\tTURNAROUND TIME\n");
for(i=0;i<n;i++)
printf("\t%d \t %d \t %d \t %d \n",i+1,ct[i],wa[i],tat[i]);
getch();}
```

INPUT:

Enter the no of processes – 3

Enter Burst Time for process 1 – 24

Enter Burst Time for process 2 -- 3

Enter Burst Time for process 3 – 3

Enter the size of time slice – 3

OUTPUT:

PROCESS	BURST TIME	WAITING TIME	TURNAROUND TIME
1	24	6	30
2	3	4	7
3	3	7	10

The Average Turnaround time is – 15.666667 The

Average Waiting time is ----- 5.666667

D). PRIORITY:

AIM: To write a c program to simulate the CPU scheduling priority algorithm.

DESCRIPTION:

To calculate the average waiting time in the priority algorithm, sort the burst times according to their priorities and then calculate the average waiting time of the processes. The waiting time of each process is obtained by summing up the burst times of all the previous processes.

ALGORITHM:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: Sort the ready queue according to the priority number.

Step 5: Set the waiting of the first process as $_0$ and its burst time as its turnaround time

Step 6: Arrange the processes based on process priority

Step 7: For each process in the Ready Q calculate Step 8:

for each process in the Ready Q calculate

a) $\text{Waiting time}(n) = \text{waiting time}(n-1) + \text{Burst time}(n-1)$

b) $\text{Turnaround time}(n) = \text{waiting time}(n) + \text{Burst time}(n)$

Step 9: Calculate

c) $\text{Average waiting time} = \text{Total waiting Time} / \text{Number of process}$

d) $\text{Average Turnaround time} = \text{Total Turnaround Time} / \text{Number of process}$ Print the results in an order.

Step10: Stop

SOURCE CODE:

```
#include<stdio.h>
main()
{
int p[20],bt[20],pri[20], wt[20],tat[20],i, k, n, temp; float wtavg,
tatavg;
clrscr();
printf("Enter the number of processes --- ");
scanf("%d",&n);
for(i=0;i<n;i++){
p[i] = i;
printf("Enter the Burst Time & Priority of Process %d --- ",i); scanf("%d
%d",&bt[i], &pri[i]);
}
for(i=0;i<n;i++)
for(k=i+1;k<n;k++)
if(pri[i] > pri[k]){
temp=p[i];
p[i]=p[k];
p[k]=temp;
temp=bt[i];
bt[i]=bt[k];
bt[k]=temp;
temp=pri[i];
pri[i]=pri[k];
pri[k]=temp;
}
wtavg = wt[0] = 0;
tatavg = tat[0] = bt[0];
for(i=1;i<n;i++)
{
wt[i] = wt[i-1] + bt[i-1];
tat[i] = tat[i-1] + bt[i];

wtavg = wtavg + wt[i];
tatavg = tatavg + tat[i];
}
printf("\nPROCESS\t\tPRIORITY\tBURST TIME\tWAITING TIME\tTURNAROUND
TIME");
for(i=0;i<n;i++)
printf("\n%d \t\t %d \t\t %d \t\t %d \t\t %d ",p[i],pri[i],bt[i],wt[i],tat[i]);
printf("\nAverage Waiting Time is --- %f",wtavg/n); printf("\nAverage
Turnaround Time is --- %f",tatavg/n);
getch();}
```


EXPERIMENT.NO 2

AIM: To Write a C program to simulate producer-consumer problem using semaphores.

DESCRIPTION

Producer consumer problem is a synchronization problem. There is a fixed size buffer where the producer produces items and that is consumed by a consumer process. One solution to the producer-consumer problem uses shared memory. To allow producer and consumer processes to run concurrently, there must be available a buffer of items that can be filled by the producer and emptied by the consumer. This buffer will reside in a region of memory that is shared by the producer and consumer processes. The producer and consumer must be synchronized, so that the consumer does not try to consume an item that has not yet been produced.

PROGRAM

```
#include<stdio.>
void main()
{
    int buffer[10], bufsize, in, out, produce, consume,
    choice=0; in = 0;
    out = 0;
    bufsize = 10;
    while(choice !=3)
    {
        printf("\n1. Produce \t 2. Consume \t3. Exit");
        printf("\nEnter your choice: ");
        scanf("%d",&choice);
        switch(choice) {
            case 1: if((in+1)%bufsize==out)
                    printf("\nBuffer is Full");
                    else
                    {
                        printf("\nEnter the value: ");
                        scanf("%d", &produce);
                        buffer[in] = produce;
                        in = (in+1)%bufsize;
                    }
                    break;;
            case 2: if(in == out)
                    printf("\nBuffer is Empty");
                    else
                    {
                        consume = buffer[out];
                        printf("\nThe consumed value is %d", consume);
                        out = (out+1)%bufsize;
                    }
                    break;
        }
    }
}
```

OUTPUT

1. Produce 2. Consume 3. Exit

Enter your choice: 2

Buffer is Empty

1. Produce 2. Consume 3. Exit

Enter your choice: 1

Enter the value: 100

1. Produce 2. Consume 3. Exit

Enter your choice: 2

The consumed value is 100

1. Produce 2. Consume 3. Exit

Enter your choice: 3

EXPERIMENT.NO 3

AIM: To Write a C program to simulate the concept of Dining-Philosophers problem.

DESCRIPTION

The dining-philosophers problem is considered a classic synchronization problem because it is an example of a large class of concurrency-control problems. It is a simple representation of the need to allocate several resources among several processes in a deadlock-free and starvation-free manner. Consider five philosophers who spend their lives thinking and eating. The philosophers share a circular table surrounded by five chairs, each belonging to one philosopher. In the center of the table is a bowl of rice, and the table is laid with five single chopsticks. When a philosopher thinks, she does not interact with her colleagues. From time to time, a philosopher gets hungry and tries to pick up the two chopsticks that are closest to her (the chopsticks that are between her and her left and right neighbors). A philosopher may pick up only one chopstick at a time. Obviously, she cannot pick up a chopstick that is already in the hand of a neighbor. When a hungry philosopher has both her chopsticks at the same time, she eats without releasing her chopsticks. When she is finished eating, she puts down both of her chopsticks and starts thinking again. The dining-philosophers problem may lead to a deadlock situation and hence some rules have to be framed to avoid the occurrence of deadlock.

PROGRAM

```
int tph, philname[20], status[20], howhung, hu[20], cho; main()
{
    int i; clrscr();
    printf("\n\nDINING PHILOSOPHER PROBLEM");
    printf("\nEnter the total no. of philosophers: ");
    scanf("%d",&tph);
    for(i=0;i<tph;i++)
    {
        philname[i]=(i+1); status[i]=1;
    }
    printf("How many are hungry : ");
    scanf("%d", &howhung);
    if(howhung==tph)
    {
        printf("\n All are hungry..\nDead lock stage will occur");
        printf("\nExiting\n");
    }
    else{
        for(i=0;i<howhung;i++){
            printf("Enterphilosopher%dposition:",(i+1));
            scanf("%d",&hu[i]);
            status[hu[i]]=2;
        }
    }
}
```

```

do
    {
        printf("1.One can eat at a time\t2.Two can eat at a time
\t3.Exit\nEnter your choice:");
        scanf("%d", &cho);
        switch(cho)
        {
            case 1: one();
                    break;
            case 2: two();
                    break;
            case 3: exit(0);
                    default: printf("\nInvalid option..");
        }
    }while(1);
}
}
one()
{
    int pos=0, x, i;
    printf("\nAllow one philosopher to eat at any time\n");
    for(i=0;i<howhung; i++, pos++)
    {
        printf("\nP %d is granted to eat", philname[hu[pos]]);
        for(x=pos;x<howhung;x++)
        printf("\nP %d is waiting", philname[hu[x]]);

    }
}
two()
{
    int i, j, s=0, t, r, x;
    printf("\n Allow two philosophers to eat at same
time\n"); for(i=0;i<howhung;i++)
    {
        for(j=i+1;j<howhung;j++)
        {
            if(abs(hu[i]-hu[j])>=1&& abs(hu[i]-hu[j])!=4)
            {
                printf("\n\ncombination %d \n", (s+1));
                t=hu[i];
                r=hu[j]; s++;
                printf("\nP %d and P %d are granted to eat", philname[hu[i]],
                    philname[hu[j]]);
            }
        }
    }
}

```

```

        for(x=0;x<howhung;x++)
        {
            if((hu[x]!=t)&&(hu[x]!=r))
                printf("\nP %d is waiting", philname[hu[x]]);
        }
    }
}

```

INPUT

DINING PHILOSOPHER PROBLEM
 Enter the total no. of philosophers: 5
 How many are hungry : 3
 Enter philosopher 1 position: 2
 Enter philosopher 2 position: 4
 Enter philosopher 3 position: 5

OUTPUT

1.One can eat at a time 2.Two can
 eat at a time 3.Exit Enter your choice: 1

Allow one philosopher to eat at any time
 P 3 is granted to eat
 P 3 is waiting
 P 5 is waiting
 P 0 is waiting
 P 5 is granted to eat
 P 5 is waiting
 P 0 is waiting
 P 0 is granted to eat
 P 0 is waiting

1. One can eat at a time 2. Two can eat at a time 3. Exit
Enter your choice: 2

Allow two philosophers to eat at same time

combination 1

P 3 and P 5 are granted to eat

P 0 is waiting

combination 2

P 3 and P 0 are granted to eat

P 5 is waiting

combination 3

P 5 and P 0 are granted to eat

P 3 is waiting

1. One can eat at a time 2. Two can
eat at a time 3. Exit Enter your choice: 3

EXPERIMENT.NO 4

MEMORY MANAGEMENT

A). MEMORY MANAGEMENT WITH FIXED PARTITIONING TECHNIQUE (MFT)

AIM: To implement and simulate the MFT algorithm.

DESCRIPTION:

In this the memory is divided in two parts and process is fit into it. The process which is best suited will be placed in the particular memory where it suits. In MFT, the memory is partitioned into fixed size partitions and each job is assigned to a partition. The memory assigned to a partition does not change. In MVT, each job gets just the amount of memory it needs. That is, the partitioning of memory is dynamic and changes as jobs enter and leave the system. MVT is a more "efficient" user of resources. MFT suffers with the problem of internal fragmentation and MVT suffers with external fragmentation.

ALGORITHM:

Step1: Start the process.

Step2: Declare variables.

Step3: Enter total memory size ms.

Step4: Allocate memory for os.

$M_s = ms - os$

Step5: Read the no partition to be divided n Partition size= ms/n .

Step6: Read the process no and process size.

Step 7: If process size is less than partition size allot else block the process. While allocating update memory wastage-external fragmentation.

```
if(pn[i]==pn[j])f=1;
```

```
if(f==0){ if(ps[i]<=siz)
```

```
{
```

```
extft=extft+size-
```

```
ps[i];avail[i]=1; count++;
```

```
}
```

```
}
```

Step 8: Print the results

SOURCE CODE :

```
#include<stdio.h>
#include<conio.h>
main()
{
int  ms,  bs,  nob,  ef,n,
mp[10],tif=0; int i,p=0;
clrscr();
printf("Enter the total memory available (in Bytes) -- ");
scanf("%d",&ms);
printf("Enter the block size (in Bytes) -- ");
scanf("%d", &bs);
nob=ms/bs;
ef=ms - nob*bs;
printf("\nEnter the number of processes -- ");
scanf("%d",&n);
for(i=0;i<n;i++)
{
printf("Enter memory required for process %d (in Bytes)-- ",i+1);
scanf("%d",&mp[i]);
}
printf("\nNo.      of      Blocks      available      in      memory--%d",nob);
printf("\n\nPROCESS\tMEMORYREQUIRED\tALLOCATED\tINTERNAL
FRAGMENTATION");
for(i=0;i<n && p<nob;i++)
{
printf("\n %d\t\t%d",i+1,mp[i]);
if(mp[i] > bs)
printf("\t\tNO\t\t---");
else
{
printf("\t\tYES\t\t%d",bs-mp[i]);
tif = tif + bs-mp[i];
p++;
}
}
if(i<n)
printf("\nMemory is Full, Remaining Processes cannot be accomodated");
printf("\n\nTotal Internal Fragmentation is %d",tif);
printf("\nTotal External Fragmentation is %d",ef);
getch();
}
```

INPUT

Enter the total memory available (in Bytes) -- 1000
Enter the block size (in Bytes)-- 300
Enter the number of processes – 5
Enter memory required for process 1 (in Bytes) -- 275
Enter memory required for process 2 (in Bytes) -- 400
Enter memory required for process 3 (in Bytes) -- 290
Enter memory required for process 4 (in Bytes) -- 293
Enter memory required for process 5 (in Bytes) -- 100
No. of Blocks available in memory -- 3

OUTPUT

PROCESS	MEMORY REQUIRED	ALLOCATED	INTERNAL FRAGMENTATION
1	275	YES	25
2	400	NO	-----
3	290	YES	10
4	293	YES	7

Memory is Full, Remaining Processes cannot be accommodated Total
Internal Fragmentation is 42
Total External Fragmentation is 100

B) MEMORY VARIABLE PARTIONING TYPE (MVT)

AIM: To write a program to simulate the MVT algorithm

ALGORITHM:

Step1: start the process.

Step2: Declare variables.

Step3: Enter total memory size ms.

Step4: Allocate memory for os.

Ms=ms-os

Step5: Read the no partition to be divided n Partition size=ms/n.

Step6: Read the process no and process size.

Step 7: If process size is less than partition size allot else block the process. While allocating update memory wastage-external fragmentation.

```
if(pn[i]==pn[j])    f=1;
```

```
if(f==0){ if(ps[i]<=size)
```

```
{
```

```
extft=extft+size-
```

```
ps[i];avail[i]=1; count++;
```

```
}
```

```
}
```

Step 8: Print the results

Step 9: Stop the process.

OUTPUT:

Enter the total memory available (in Bytes) – 1000
Enter memory required for process 1 (in Bytes) – 400
Memory is allocated for Process 1
Do you want to continue(y/n) -- y
Enter memory required for process 2 (in Bytes) -- 275
Memory is allocated for Process 2
Do you want to continue(y/n) – y
Enter memory required for process 3 (in Bytes) – 550

Memory is Full

Total Memory Available – 1000

PROCESS	MEMORY ALLOCATED
1	400
2	275

Total Memory Allocated is 675

Total External Fragmentation is 325

VIVA QUESTIONS

- 1) What is MFT?
- 2) What is MVT?
- 3) What is the difference between MVT and MFT?
- 4) What is meant by fragmentation?
- 5) Give the difference between internal and external fragmentation

EXPERIMENT.NO 5

MEMORY ALLOCATION TECHNIQUES

AIM: To Write a C program to simulate the following contiguous memory allocation techniques

- a) Worst-fit b) Best-fit c) First-fit

DESCRIPTION

One of the simplest methods for memory allocation is to divide memory into several fixed-sized partitions. Each partition may contain exactly one process. In this multiple-partition method, when a partition is free, a process is selected from the input queue and is loaded into the free partition. When the process terminates, the partition becomes available for another process. The operating system keeps a table indicating which parts of memory are available and which are occupied. Finally, when a process arrives and needs memory, a memory section large enough for this process is provided. When it is time to load or swap a process into main memory, and if there is more than one free block of memory of sufficient size, then the operating system must decide which free block to allocate. Best-fit strategy chooses the block that is closest in size to the request. First-fit chooses the first available block that is large enough. Worst-fit chooses the largest available block.

PROGRAM

WORST-FIT

```
#include<stdio.h>
#include<conio.h>
#define max 25
void main()
{
    int
    frag[max],b[max],f[max],i,j,nb,nf,t
    emp; static int bf[max],ff[max];
    clrscr();
    printf("\n\tMemory Management Scheme - First Fit");
    printf("\n\tEnter the number of blocks:");
    scanf("%d",&nb);
    printf("Enter the number of files:");
    scanf("%d",&nf);
    printf("\n\tEnter the size of the blocks:-\n");
    for(i=1;i<=nb;i++)
    {
        printf("Block %d:",i);
        scanf("%d",&b[i]);
    }
    printf("Enter the size of the files :-\n");
    for(i=1;i<=nf;i++)
    {
        printf("File %d:",i);
        scanf("%d",&f[i]);
```

```

    }
    for(i=1;i<=nf;i++)
    {
        for(j=1;j<=nb;j++)
        {
            if(bf[j]!=1)
            {
                temp=b[j]-f[i];
                if(temp>=0)
                {
                    ff[i]=j;
                    break;
                }
            }
        }
        frag[i]=temp;
        bf[ff[i]]=1;
    }
    printf("\nFile_no:\tFile_size :\tBlock_no:\tBlock_size:\tFragement");
    for(i=1;i<=nf;i++)
    printf("\n%d\t%d\t%d\t%d\t%d",i,f[i],ff[i],b[ff[i]],frag[i]);
    getch();
}

```

INPUT

Enter the number of blocks: 3

Enter the number of files: 2

Enter the size of the blocks:-

Block 1: 5

Block 2: 2

Block 3: 7

Enter the size of the files:-

File 1: 1

File 2: 4

OUTPUT

File No	File Size	Block No	Block Size	Fragment
1	1	1	5	4
2	4	3	7	3

BEST-FIT

```
#include<stdio.h>
#include<conio.h>
#define max 25
void main()
{
    int frag[max],b[max],f[max],i,j,nb,nf,temp,lowest=10000;
    static int bf[max],ff[max];
    clrscr();
    printf("\nEnter the number of blocks:");
    scanf("%d",&nb);
    printf("Enter the number of files:");
    scanf("%d",&nf);
    printf("\nEnter the size of the blocks:-\n");
    for(i=1;i<=nb;i++)
    printf("Block %d:",i);
    scanf("%d",&b[i]);
    printf("Enter the size of the files :-\n");
    for(i=1;i<=nf;i++)
    {
        printf("File %d:",i);
        scanf("%d",&f[i]);
    }
    for(i=1;i<=nf;i++)
    {
        for(j=1;j<=nb;j++)
        {
            if(bf[j]!=1)
            {
                temp=b[j]-f[i];
                if(temp>=0)
                if(lowest>temp)
                {
                    ff[i]=j;
                    lowest=temp;
                }
            }
        }
        frag[i]=lowest; bf[ff[i]]=1; lowest=10000;
    }
    printf("\nFile No\tFile Size \tBlock No\tBlock
    Size\tFragment"); for(i=1;i<=nf && ff[i]!=0;i++)

        printf("\n%d\t%d\t%d\t%d\t%d",i,f[i],ff[i],b[ff[i]],frag[i]);
    getch();
}
```

INPUT

Enter the number of blocks: 3

Enter the number of files: 2

Enter the size of the blocks:-

Block 1: 5

Block 2: 2

Block 3: 7

Enter the size of the files:-

File 1: 1

File 2: 4

OUTPUT

	File No	File Size	Block No	Block Size	Fragment
1	1		2	2	1
2	4		1	5	1

FIRST-FIT

```
#include<stdio.h>
#include<conio.h>
#define max 25
void main()
{
    int
    frag[max],b[max],f[max],i,j,nb,nf,temp,highes
    t=0; static int bf[max],ff[max];
    clrscr();
    printf("\n\tMemory Management Scheme - Worst Fit");
    printf("\nEnter the number of blocks:");
    scanf("%d",&nb);
    printf("Enter the number of files:");
    scanf("%d",&nf);
    printf("\nEnter the size of the blocks:-\n");
    for(i=1;i<=nb;i++)
    {
        printf("Block %d:",i);
        scanf("%d",&b[i]);
    }
    printf("Enter the size of the files :-\n");
    for(i=1;i<=nf;i++)
    {
        printf("File %d:",i);
        scanf("%d",&f[i]);
    }
}
```

```

for(i=1;i<=nf;i++)
{
    for(j=1;j<=nb;j++)
    {
        if(bf[j]!=1) //if bf[j] is not allocated
        {
            temp=b[j]-f[i];
            if(temp>=0)
                if(highest<temp)
                    {
                }
            }
        }
        frag[i]=highest; bf[ff[i]]=1; highest=0;
    }
    ff[i]=j; highest=temp;
}
printf("\nFile_no:\tFile_size:\tBlock_no:\tBlock_size:\tFragement");
for(i=1;i<=nf;i++)
    printf("\n%d\t%d\t%d\t%d\t%d",i,f[i],ff[i],b[ff[i]],frag[i]);
getch();
}

```

INPUT

Enter the number of blocks: 3

Enter the number of files: 2

Enter the size of the blocks:-

Block 1: 5

Block 2: 2

Block 3: 7

Enter the size of the files:-

File 1: 1

File 2: 4

OUTPUT

File No	File Size	Block No	Block Size	Fragment
1	1	3	7	6
2	4	1	5	1

EXPERIMENT NO.6

PAGE REPLACEMENT ALGORITHMS

AIM: To implement FIFO page replacement technique.

a) FIFO b) LRU c) OPTIMAL

DESCRIPTION:

Page replacement algorithms are an important part of virtual memory management and it helps the OS to decide which memory page can be moved out making space for the currently needed page. However, the ultimate objective of all page replacement algorithms is to reduce the number of page faults.

FIFO-This is the simplest page replacement algorithm. In this algorithm, the operating system keeps track of all pages in the memory in a queue, the oldest page is in the front of the queue. When a page needs to be replaced page in the front of the queue is selected for removal.

LRU-In this algorithm page will be replaced which is least recently used

OPTIMAL- In this algorithm, pages are replaced which would not be used for the longest duration of time in the future. This algorithm will give us less page fault when compared to other page replacement algorithms.

ALGORITHM:

- 1.** Start the process
- 2.** Read number of pages n
- 3.** Read number of pages no
- 4.** Read page numbers into an array a[i]
- 5.** Initialize avail[i]=0 .to check page hit
- 6.** Replace the page with circular queue, while re-placing check page availability in the frame
Place avail[i]=1 if page is placed in the frame Count page faults
- 7.** Print the results.
- 8.** Stop the process.

A) FIRST IN FIRST OUT
SOURCE CODE :

```
#include<stdio.h>
#include<conio.h> int fr[3];
void main()
{
void display();
int i,j,page[12]={2,3,2,1,5,2,4,5,3,2,5,2};
int
flag1=0,flag2=0,pf=0,frsize=3,top=0;
clrscr();
for(i=0;i<3;i++)
{
fr[i]=-1;
}
for(j=0;j<12;j++)
{
flag1=0; flag2=0; for(i=0;i<12;i++)
{
if(fr[i]==page[j])
{
flag1=1; flag2=1; break;
}
}
if(flag1==0)
{
for(i=0;i<frsize;i++)
{
if(fr[i]==-1)
{
fr[i]=page[j]; flag2=1; break;
}
}
}
if(flag2==0)
{
fr[top]=page[j];
top++;
pf++;
if(top>=frsize)
top=0;
}
display();
}
```

```
printf("Number of page faults : %d ",pf+frsize);
getch();
}
void display()
{
int i; printf("\n");
for(i=0;i<3;i++)
printf("%d\t",fr[i]);
}
```

OUTPUT:

```
2 -1 -1
2 3 -1
2 3 -1
2 3 1
5 3 1
5 2 1
5 2 4
5 2 4
3 2 4
3 2 4
3 5 4
3 5 2
```

Number of page faults: 9

B) LEAST RECENTLY USED

AIM: To implement LRU page replacement technique.

ALGORITHM:

1. Start the process
2. Declare the size
3. Get the number of pages to be inserted
4. Get the value
5. Declare counter and stack
6. Select the least recently used page by counter value
7. Stack them according the selection.
8. Display the values
9. Stop the process

SOURCE CODE :

```
#include<stdio.h>
#include<conio.h>
int fr[3];
void main()
{
void display();
int p[12]={2,3,2,1,5,2,4,5,3,2,5,2},i,j,fs[3];
int index,k,l,flag1=0,flag2=0,pf=0,frsize=3;
clrscr();
for(i=0;i<3;i++)
{
fr[i]=-1;
}
for(j=0;j<12;j++)
{
flag1=0,flag2=0;
for(i=0;i<3;i++)
{
if(fr[i]==p[j])
{
flag1=1;
flag2=1; break;
}
}
if(flag1==0)
```

```

{
for(i=0;i<3;i++)
{
if(fr[i]==-1)
{
fr[i]=p[j];    flag2=1;
break;
}
}
}
if(flag2==0)
{
for(i=0;i<3;i++)
fs[i]=0;
for(k=j-1,l=1;l<=frsize-1;l++,k--)
{
for(i=0;i<3;i++)
{
if(fr[i]==p[k]) fs[i]=1;
}}
for(i=0;i<3;i++)
{
if(fs[i]==0)
index=i;
}
fr[index]=p[j];
pf++;
}
display();
}
printf("\n no of page faults :%d",pf+frsize);
getch();
}
void display()
{
int i; printf("\n");
for(i=0;i<3;i++)
printf("\t%d",fr[i]);
}

```

OUTPUT:

2 -1 -1
2 3 -1
2 3 -1
2 3 1
2 5 1
2 5 1
2 5 4
2 5 4
3 5 4
3 5 2
3 5 2
3 5 2

No of page faults: 7

C) OPTIMAL

AIM: To implement optimal page replacement technique.

ALGORITHM:

1. Start Program
2. Read Number Of Pages And Frames
3. Read Each Page Value
4. Search For Page In The Frames
5. If Not Available Allocate Free Frame
6. If No Frames Is Free Replace The Page With The Page That Is Leastly Used
7. Print Page Number Of Page Faults
8. Stop process.

SOURCE CODE:

```
/* Program to simulate optimal page replacement */
#include<stdio.h>
#include<conio.h>
int fr[3], n, m;
void
display();
void main()
{
int i,j,page[20],fs[10];
int
max,found=0,l[3],index,k,l,flag1=0,flag2=0,pf=0;
float pr;
clrscr();
printf("Enter length of the reference string: ");
scanf("%d",&n);
printf("Enter the reference string: ");
for(i=0;i<n;i++)
scanf("%d",&page[i]);
printf("Enter no of frames: ");
scanf("%d",&m);
for(i=0;i<m;i++)
fr[i]=-1; pf=m;
```

```

for(j=0;j<n;j++)
{
flag1=0;    flag2=0;
for(i=0;i<m;i++)
{
if(fr[i]==page[j])
{
flag1=1; flag2=1;
break;
}
}
if(flag1==0)
{
for(i=0;i<m;i++)
{
if(fr[i]==-1)
{
fr[i]=page[j]; flag2=1;
break;
}
}
}
if(flag2==0)
{
for(i=0;i<m;i++)
lg[i]=0;
for(i=0;i<m;i++)
{
for(k=j+1;k<=n;k++)
{
if(fr[i]==page[k])
{
lg[i]=k-j;
break;
}
}
}
}
found=0;
for(i=0;i<m;i++)
{
if(lg[i]==0)
{
index=i;
found = 1;

```

```

break;
}
}
if(found==0)
{
max=lg[0]; index=0;
for(i=0;i<m;i++)
{
if(max<lg[i])
{
max=lg[i];
index=i;
}
}
}
fr[index]=page[j];
pf++;
}
display();
}
printf("Number of page faults : %d\n", pf);
pr=(float)pf/n*100;
printf("Page fault rate = %f\n", pr); getch();
}
void display()
{
int i; for(i=0;i<m;i++)
printf("%d\t",fr[i]);
printf("\n");
}

```

OUTPUT:

Enter length of the reference string: 12

Enter the reference string: 1 2 3 4 1 2 5 1 2 3 4 5

Enter no of frames: 3

1 -1 -1

1 2 -1

1 2 3

1 2 4

1 2 4

1 2 4

1 2 5

1 2 5

1 2 5

3 2 5

4 2 5

4 2 5

Number of page faults : 7 Page fault rate = 58.333332

VIVA QUESTIONS

- 1) What is meant by page fault?
- 2) What is meant by paging?
- 3) What is page hit and page fault rate?
- 4) List the various page replacement algorithm
- 5) Which one is the best replacement algorithm?

EXPERIMENT NO. 7

FILE ORGANIZATION TECHNIQUES

A) SINGLE LEVEL DIRECTORY:

AIM: Program to simulate Single level directory file organization technique.

DESCRIPTION:

The directory structure is the organization of files into a hierarchy of folders. In a single-level directory system, all the files are placed in one directory. There is a root directory which has all files. It has a simple architecture and there are no sub directories. Advantage of single level directory system is that it is easy to find a file in the directory.

SOURCE CODE :

```
#include<stdio.h>
struct
{
char  dname[10],fname[10][10];
int fcnt;
}dir;

void main()
{
int i,ch; char
f[30]; clrscr();
dir.fcnt = 0;
printf("\nEnter name of directory -- ");
scanf("%s", dir.dname);
while(1)
{
printf("\n\n1. Create File\t2. Delete File\t3. Search File \n
4. Display Files\t5. Exit\nEnter your choice -- ");
scanf("%d",&ch);
switch(ch)
{
case 1: printf("\nEnter the name of the file -- ");
scanf("%s",dir.fname[dir.fcnt]);
dir.fcnt++; break;
case 2: printf("\nEnter the name of the file -- ");
scanf("%s",f);
for(i=0;i<dir.fcnt;i++)
{
if(strcmp(f, dir.fname[i])==0)
{
printf("File %s is deleted ",f); strcpy(dir.fname[i],dir.fname[dir.fcnt-1]); break;
}
}
```

```

}
if(i==dir.fcnt)
printf("File %s not found",f);
else
    dir.fcnt--;
    break;
    case 3:
        printf("\nEnter the name of the file -- ");
        scanf("%s",f);
        for(i=0;i<dir.fcnt;i++)
        {
            if(strcmp(f, dir.fname[i])==0)
            {
                printf("File %s is found ", f);
                break;
            }
        }
        if(i==dir.fcnt)
        printf("File %s not found",f);
        break;
    case 4:
        if(dir.fcnt==0)
        printf("\nDirectory Empty");
        else
        {
            printf("\nThe Files are -- ");
            for(i=0;i<dir.fcnt;i++)
            printf("\t%s",dir.fname[i]);
        }
        break;
    default: exit(0);
}
}
getch();}

```

OUTPUT:

Enter name of directory -- CSE

1. Create File
 2. Delete File
 3. Search File
 4. Display Files
 5. Exit
- Enter your choice – 1

Enter the name of the file -- A

1. Create File
 2. Delete File
 3. Search File
 4. Display Files
 5. Exit
- Enter your choice – 1

Enter the name of the file -- B

1. Create File
 2. Delete File
 3. Search File
 4. Display Files
 5. Exit
- Enter your choice – 1

Enter the name of the file -- C

1. Create File
 2. Delete File
 3. Search File
 4. Display Files
 5. Exit
- Enter your choice – 4

The Files are -- A B C

1. Create File
 2. Delete File
 3. Search File
 4. Display Files
 5. Exit
- Enter your choice – 3

Enter the name of the file – ABC File

ABC not found

1. Create File
 2. Delete File
 3. Search File
 4. Display Files
 5. Exit
- Enter your choice – 2

Enter the name of the file – B

File B is deleted

1. Create File
 2. Delete File
 3. Search File
 4. Display Files
 5. Exit
- Enter your choice – 5


```

        dir[i].fcnt++;
        printf("File created");
    }
    if(i==dcnt)
        printf("Directory %s not found",d);
        break;
case 3: printf("\nEnter name of the directory -- ");
        scanf("%s",d);
        for(i=0;i<dcnt;i++)
        for(i=0;i<dcnt;i++)
        {
            if(strcmp(d,dir[i].dname)==0)
            {
                printf("Enter name of the file -- ");
                scanf("%s",f);
                for(k=0;k<dir[i].fcnt;k++)
                {
                    if(strcmp(f, dir[i].fname[k])==0)
                    {
                        printf("File %s is deleted ",f);
                        dir[i].fcnt--;
                        strcpy(dir[i].fname[k],dir[i].fname[dir[i].fcnt]);
                        goto jmp;
                    }
                }

                printf("File %s not found",f); goto jmp;
            }
        }
        printf("Directory %s not found",d);
        jmp : break;
case 4: printf("\nEnter name of the directory -- ");
        scanf("%s",d);
        for(i=0;i<dcnt;i++)
        {
            if(strcmp(d,dir[i].dname)==0)
            {
                printf("Enter the name of the file -- ");
                scanf("%s",f);
                for(k=0;k<dir[i].fcnt;k++)
                {
                    if(strcmp(f, dir[i].fname[k])==0)
                    {
                        printf("File %s is found ",f); goto jmp1;
                    }
                }

                printf("File %s not found",f); goto jmp1;
            }
        }
}
}

```

```

        printf("Directory %s not found",d); jmp1: break;
    case 5: if(dcnt==0)
        printf("\nNo Directory's ");
        else
        {
            printf("\nDirectory\Files");
            for(i=0;i<dcnt;i++)
            {
                printf("\n%s\t",dir[i].dname);
                for(k=0;k<dir[i].fcnt;k++)
                printf("\t%s",dir[i].fname[k]);
            }
        }
        break;

    default:exit(0);
}
}
getch();
}

```

OUTPUT

1. Create Directory 2. Create File 3. Delete File
4. Search File 5. Display 6. Exit
Enter your choice -- 1
Enter name of directory -- DIR1 Directory created

1. Create Directory 2. Create File 3. Delete File
4. Search File 5. Display 6. Exit Enter your choice -- 1
Enter name of directory -- DIR2 Directory created

1. Create Directory 2. Create File 3. Delete File
4. Search File 5. Display 6. Exit Enter your choice -- 2
Enter name of the directory -- DIR1
Enter name of the file -- A1
File created

1. Create Directory 2. Create File 3. Delete File
4. Search File 5. Display 6. Exit
Enter your choice -- 2
Enter name of the directory -- DIR1

Enter name of the file -- A2
File created

1. Create Directory 2. Create File 3. Delete File
4. Search File 5. Display 6.
Exit Enter your choice -- 6

VIVA QUESTIONS

1. Define directory?
2. List the different types of directory structures?
3. What is the advantage of hierarchical directory structure?
4. Which of the directory structures is efficient? Why?
5. What is acyclic graph directory?

EXPERIMENT.NO.8

FILE ALLOCATION STRATEGIES

A) SEQUENTIAL:

AIM: To write a C program for implementing sequential file allocation method

DESCRIPTION:

The most common form of file structure is the sequential file in this type of file, a fixed format is used for records. All records (of the system) have the same length, consisting of the same number of fixed length fields in a particular order because the length and position of each field are known, only the values of fields need to be stored, the field name and length for each field are attributes of the file structure.

ALGORITHM:

Step 1: Start the program.

Step 2: Get the number of files.

Step 3: Get the memory requirement of each file.

Step 4: Allocate the required locations to each in sequential order a).

Randomly select a location from available location $s1 = \text{random}(100)$;

a) Check whether the required locations are free from the selected location.

```
if(b[s1].flag==0){  
for      (j=s1;j<s1+p[i];j++){  
if((b[j].flag)==0)count++;  
}
```

```
if(count==p[i]) break;
```

```
}
```

b) Allocate and set flag=1 to the allocated locations. $\text{for}(s=s1;s<(s1+p[i]);s++)$

```
{
```

```
k[i][j]=s; j=j+1; b[s].bno=s;
```

```
b[s].flag=1;
```

```
}
```

Step 5: Print the results file no, length, Blocks allocated. Step

6: Stop the program

SOURCE CODE :

```
#include<stdio.h>
main()
{
int f[50],i,st,j,len,c,k;
clrscr();
for(i=0;i<50;i++)
f[i]=0;
X:
printf("\n Enter the starting block & length of file");
scanf("%d%d",&st,&len);
for(j=st;j<(st+len);j++)
if(f[j]==0)
{
f[j]=1
;
printf("\n%d->%d",j,f[j]);
}
else
{
printf("Block already allocated");
break;
}
if(j==(st+len))
printf("\n the file is allocated to disk");
printf("\n if u want to enter more files?(y-1/n-0)");
scanf("%d",&c);
if(c==1)
goto X;
else
exit();
getch();
}
```

OUTPUT:

Enter the starting block & length of file 4 10

4->1

5->1

6->1

7->1

8->1

9->1

10->1

11->1

12->1

13->1

The file is allocated to disk.

B) INDEXED:

AIM: To implement allocation method using chained method

DESCRIPTION:

In the chained method file allocation table contains a field which points to starting block of memory. From it for each bloc a pointer is kept to next successive block. Hence, there is no external fragmentation.

ALGORITHM:

Step 1: Start the program.

Step 2: Get the number of files.

Step 3: Get the memory requirement of each file.

Step 4: Allocate the required locations by selecting a location randomly $q = \text{random}(100)$;

a) Check whether the selected location is free .

b) If the location is free allocate and set $\text{flag}=1$ to the allocated locations.

```
q=random(100);
```

```
{
```

```
if(b[q].flag==0)
```

```
b[q].flag=1;
```

```
b[q].fno=j;
```

```
r[i][j]=q;
```

Step 5: Print the results file no, length ,Blocks allocated.

Step 6: Stop the program

SOURCE CODE :

```
#include<stdio.h>
int f[50],i,k,j,inde[50],n,c,count=0,p;
main()
{
clrscr();
for(i=0;i<50;i++)
f[i]=0;
x: printf("enter index block\t");
scanf("%d",&p);
if(f[p]==0)
{
f[p]=1;
printf("enter no of files on index\t");
scanf("%d",&n);
}
else
{
printf("Block already allocated\n");
goto x;
}
for(i=0;i<n;i++)
scanf("%d",&inde[i]);
for(i=0;i<n;i++)
if(f[inde[i]]==1)
{
printf("Block already allocated");
goto x;
}
for(j=0;j<n;j++)
f[inde[j]]=1;
printf("\n allocated");
printf("\n file indexed");
for(k=0;k<n;k++)
printf("\n %d->%d:%d",p,inde[k],f[inde[k]]);
printf(" Enter 1 to enter more files and 0 to exit\t");
scanf("%d",&c);
if(c==1)
goto x;
else
exit();
getch();
}
```

OUTPUT: enter index block 9

Enter no of files on index 3 1

2 3

Allocated

File indexed

9->1:1

9->2;1

9->3:1 enter 1 to enter more files and 0 to exit

C) LINKED:

AIM: To implement linked file allocation technique.

DESCRIPTION:

In the chained method file allocation table contains a field which points to starting block of memory. From it for each bloc a pointer is kept to next successive block. Hence, there is no external fragmentation

ALGORITHM:

Step 1: Start the program.

Step 2: Get the number of files.

Step 3: Get the memory requirement of each file.

Step 4: Allocate the required locations by selecting a location randomly $q = \text{random}(100)$;

a) Check whether the selected location is free .

b) If the location is free allocate and set $\text{flag}=1$ to the allocated locations.

While allocating next location address to attach it to previous location

```
for(i=0;i<n;i++)
{
for(j=0;j<s[i];j++)
{
q=random(100);          if(b[q].flag==0)
b[q].flag=1;
b[q].fno=j;
r[i][j]=q;
        if(j>0)
        {
}
}
}
```

```
p=r[i][j-1]; b[p].next=q;}
```

Step 5: Print the results file no, length ,Blocks allocated.

Step 6: Stop the program

SOURCE CODE :

```
#include<stdio.h>
main()
{
int f[50],p,i,j,k,a,st,len,n,c;
clrscr();
for(i=0;i<50;i++) f[i]=0;
printf("Enter how many blocks that are already
allocated"); scanf("%d",&p);
printf("\nEnter the blocks no.s that are already allocated");
for(i=0;i<p;i++)
{
scanf("%d",&a);
f[a]=1;
}
X:
printf("Enter the starting index block &
length"); scanf("%d%d",&st,&len); k=len;
for(j=st;j<(k+st);j++)
{
if(f[j]==0)
{ f[j]=1;
printf("\n%d->%d",j,f[j]);
}
else
{
printf("\n %d->file is already
allocated",j);
k++;
}
}
printf("\n If u want to enter one
more file? (yes-1/no-0)");
scanf("%d",&c);
if(c==1)
goto
X;
else
exit();
getch();}
```

OUTPUT:

Enter how many blocks that are already allocated 3 Enter the blocks no.s that are already allocated 4 7 Enter the starting index block & length 3 7 9

3->1

4->1 file is already allocated

5->1

6->1

7->1 file is already allocated

8->1

9->1file is already allocated

10->1

11->1

12->1

VIVA QUESTIONS

- 1) List the various types of files
- 2) What are the various file allocation strategies?
- 3) What is linked allocation?
- 4) What are the advantages of linked allocation?
- 5) What are the disadvantages of sequential allocation methods?

EXPERIMENT.NO 9

DEAD LOCK AVOIDANCE

AIM: To Simulate bankers algorithm for Dead Lock Avoidance (Banker's Algorithm)

DESCRIPTION:

Deadlock is a situation where in two or more competing actions are waiting for the other to finish, and thus neither ever does. When a new process enters a system, it must declare the maximum number of instances of each resource type it needed. This number may exceed the total number of resources in the system. When the user request a set of resources, the system must determine whether the allocation of each resources will leave the system in safe state. If it will the resources are allocation; otherwise the process must wait until some other process release the resources.

Data structures

n-Number of process, m-number of resource types.

Available: Available[j]=k, k – instance of resource type R_j is available. Max: If

max[i, j]=k, P_i may request at most k instances resource R_j.

Allocation: If Allocation [i, j]=k, P_i allocated to k instances of resource R_j Need: If

Need[I, j]=k, P_i may need k more instances of resource type R_j, Need[I, j]=Max[I, j]-

Allocation[I, j];

Safety Algorithm

1. Work and Finish be the vector of length m and n respectively, Work=Available and Finish[i] =False.
2. Find an i such that both
Finish[i] =False
Need<=Work If no such I exists go to step 4.
3. work= work + Allocation, Finish[i] =True;
4. if Finish[1]=True for all I, then the system is in safe state.

Resource request algorithm

Let Request i be request vector for the process P_i, If request i=[j]=k, then process P_i wants k instances of resource type R_j.

1. if Request<=Need I go to step 2. Otherwise raise an error condition.
2. if Request<=Available go to step 3. Otherwise P_i must since the resources are available.
3. Have the system pretend to have allocated the requested resources to process P_i by modifying the state as follows;
Available=Available-Request I;
Allocation I=Allocation +Request I;
Need i=Need i- Request I;

If the resulting resource allocation state is safe, the transaction is completed and process P_i is allocated its resources. However if the state is unsafe, the P_i must wait for Request i and the old resource-allocation state is restored.

ALGORITHM:

1. Start the program.
2. Get the values of resources and processes.
3. Get the avail value.
4. After allocation find the need value.
5. Check whether its possible to allocate.
6. If it is possible then the system is in safe state.
7. Else system is not in safety state.
8. If the new request comes then check that the system is in safety.
9. or not if we allow the request.
10. stop the program.
11. *end*

SOURCE CODE :

```
#include<stdio.h>
#include<conio.h>
#include<string.h>
void main()
{
int alloc[10][10],max[10][10];
int avail[10],work[10],total[10];
int i,j,k,n,need[10][10];
int m;
int count=0,c=0;
char finish[10];
clrscr();
printf("Enter the no. of processes and resources:");
scanf("%d%d",&n,&m);
for(i=0;i<=n;i++)
finish[i]='n';
printf("Enter the claim matrix:\n");
for(i=0;i<n;i++)
for(j=0;j<m;j++)
scanf("%d",&max[i][j]);
printf("Enter the allocation matrix:\n");
for(i=0;i<n;i++)
for(j=0;j<m;j++)
scanf("%d",&alloc[i][j]);
printf("Resource vector:");
for(i=0;i<m;i++)
scanf("%d",&total[i]);
for(i=0;i<m;i++)
avail[i]=0;
for(i=0;i<n;i++)
```

```

for(j=0;j<m;j++)
avail[j]+=alloc[i][j];
for(i=0;i<m;i++)
work[i]=avail[i];
for(j=0;j<m;j++)
work[j]=total[j]-work[j];
for(i=0;i<n;i++)
for(j=0;j<m;j++)
need[i][j]=max[i][j]-alloc[i][j];
A:
for(i=0;i<n;i++)
{
c=0;
for(j=0;j<m;j++)
if((need[i][j]<=work[j])&&(finish[i]=='n'))
c++;
if(c==m)
{
printf("All the resources can be allocated to Process %d", i+1);
printf("\n\nAvailable resources are:");
for(k=0;k<m;k++)
{
work[k]+=alloc[i][k];
printf("%4d",work[k]);
}
printf("\n");
finish[i]='y';
printf("\nProcess %d executed?:%c \n",i+1,finish[i]);
count++;
}
}
if(count!=n)
goto A;
else
printf("\n System is in safe mode");
printf("\n The given state is safe state");
getch();
}

```

OUTPUT

Enter the no. of processes and resources: 4 3

Enter the claim matrix:

3 2 2

6 1 3

3 1 4

4 2 2

Enter the allocation matrix:

1 0 0

6 1 2

2 1 1

0 0 2

Resource vector:9 3 6

All the resources can be allocated to Process 2

Available resources are: 6 2 3

Process 2 executed?:y

All the resources can be allocated to Process 3 Available resources are: 8 3 4

Process 3 executed?:y

All the resources can be allocated to Process 4 Available resources are: 8 3 6

Process 4 executed?:y

All the resources can be allocated to Process 1

Available resources are: 9 3 6

Process 1 executed?:y

System is in safe mode

The given state is safe state

VIVA QUESTIONS

- 1) What is meant by deadlock?
- 2) What is safe state in banker's algorithms?
- 3) What is banker's algorithm?
- 4) What are the necessary conditions where deadlock occurs?
- 5) What are the principles and goals of protection?

EXPERIMENT.NO 10

DEAD LOCK PREVENTION

AIM: To implement deadlock prevention technique

Banker's Algorithm:

When a new process enters a system, it must declare the maximum number of instances of each resource type it needed. This number may exceed the total number of resources in the system. When the user request a set of resources, the system must determine whether the allocation of each resources will leave the system in safe state. If it will the resources are allocation; otherwise the process must wait until some other process release the resources.

DESCRIPTION:

Data structures

-
- n-Number of process, m-number of resource types.
-
- Available: Available[j]=k, k – instance of resource type R_j is available.
- Max: If max[i, j]=k, P_i may request at most k instances resource R_j.
Allocation: If Allocation [i, j]=k, P_i allocated to k instances of resource R_j Need:
If Need[I, j]=k, P_i may need k more instances of resource type R_j,
Need[I, j]=Max[I, j]-Allocation[I, j];

Safety Algorithm

Work and Finish be the vector of length m and n respectively, Work=Available and Finish[i] =False.

Find an i such
that both •Finish[i] =False
Need<=Work

If no such I exists go to step 4.

5. work=work+Allocation, Finish[i] =True;

if Finish[1]=True for all I, then the system is in safe state

ALGORITHM:

1. Start the program.
2. Get the values of resources and processes.
3. Get the avail value.
4. After allocation find the need value.
5. Check whether its possible to allocate.
6. If it is possible then the system is in safe state.
7. Else system is not in safety state
8. Stop the process.

SOURCE CODE :

```
#include<stdio.h>
#include<conio.h>
void main()
{
char job[10][10];
int time[10],avail,tem[10],temp[10]; int safe[10];
int ind=1,i,j,q,n,t;
clrscr();
printf("Enter no of jobs: ");
scanf("%d",&n);
for(i=0;i<n;i++)
{
printf("Enter name and time: ");
scanf("%s%d",&job[i],&time[i]);
}
printf("Enter the available resources:");
scanf("%d",&avail);
for(i=0;i<n;i++)
{
temp[i]=time[i];
tem[i]=i;
}
for(i=0;i<n;i++)
for(j=i+1;j<n;j++)
{
if(temp[i]>temp[j])
{
t=temp[i];
```

```

temp[i]=temp[j];
temp[j]=t; t=tem[i];
tem[i]=tem[j];
tem[j]=t;
}
}
for(i=0;i<n;i++)
{
q=tem[i];
if(time[q]<=avail)
{
safe[ind]=tem[i];
avail=avail-tem[q];
printf("%s",job[safe[ind]]);
ind++;
}
else
{
printf("No safe sequence\n");
}
}
printf("Safe sequence is:");
for(i=1;i<ind; i++)
printf("%s %d\n",job[safe[i]],time[safe[i]]);
getch();
}

```

OUTPUT:

```

Enter no of jobs:4
Enter name and time: A 1
Enter name and time: B 4
Enter name and time: C 2
Enter name and time: D 3
Enter the available resources: 20
Safe sequence is: A 1, C 2, D 3, B 4.

```

EXPERIMENT.NO 11

AIM: To Write a C program to simulate disk scheduling algorithms

a) FCFS b) SCAN c) C-SCAN

DESCRIPTION

One of the responsibilities of the operating system is to use the hardware efficiently. For the disk drives, meeting this responsibility entails having fast access time and large disk bandwidth. Both the access time and the bandwidth can be improved by managing the order in which disk I/O requests are serviced which is called as disk scheduling. The simplest form of disk scheduling is, of course, the first-come, first-served (FCFS) algorithm. This algorithm is intrinsically fair, but it generally does not provide the fastest service. In the SCAN algorithm, the disk arm starts at one end, and moves towards the other end, servicing requests as it reaches each cylinder, until it gets to the other end of the disk. At the other end, the direction of head movement is reversed, and servicing continues. The head continuously scans back and forth across the disk. C-SCAN is a variant of SCAN designed to provide a more uniform wait time. Like SCAN, C-SCAN moves the head from one end of the disk to the other, servicing requests along the way. When the head reaches the other end, however, it immediately returns to the beginning of the disk without servicing any requests on the return trip

PROGRAM

A) FCFS DISK SCHEDULING ALGORITHM

```
#include<stdio.h>
main()
{
    int t[20], n, I, j, tohm[20], tot=0; float avhm;
    clrscr();
    printf("enter the no.of tracks");
    scanf("%d",&n);
    printf("enter the tracks to be traversed");
    for(i=2;i<n+2;i++)
        scanf("%d",&t*i+);
    for(i=1;i<n+1;i++)
    {
        tohm[i]=t[i+1]-t[i];
        if(tohm[i]<0)
            tohm[i]=tohm[i]*(-1);
    }
    for(i=1;i<n+1;i++)
        tot+=tohm[i];
    avhm=(float)tot/n;
    printf("Tracks traversed\tDifference between tracks\n");
    for(i=1;i<n+1;i++)
        printf("%d\t\t\t%d\n",t*i+,tohm*i+);
    printf("\nAverage header movements:%f",avhm);
    getch();
}
```

INPUT

Enter no.of tracks:9

Enter track position:55 58 60 70 18 90 150 160 184

OUTPUT

Tracks traversed

55

58

60

70

18

90

150

160

184

Difference between tracks

45

3

2

10

52

72

60

10

24

Average header movements:30.888889

B) SCAN DISK SCHEDULING ALGORITHM

```
#include<stdio.h>
main()
{
    int t[20], d[20], h, i, j, n, temp, k, atr[20], tot, p, sum=0;
    clrscr();
    printf("enter the no of tracks to be traveresed");
    scanf("%d",&n);
    printf("enter the position of head");
    scanf("%d",&h);
    t[0]=0;t[1]=h;
    printf("enter the tracks");
    for(i=2;i<n+2;i++)
        scanf("%d",&t[i]);
    for(i=0;i<n+2;i++)
    {
        for(j=0;j<(n+2)-i-1;j++)
        {
            if(t[j]>t[j+1])
            {
                temp=t[j];
                t[j]=t[j+1];
                t[j+1]=temp;
            } } }
    for(i=0;i<n+2;i++)
    if(t[i]==h)
        j=i;k=i;
    p=0;
    while(t[j]!=0)
    {
        atr[p]=t[j]; j--;
        p++;
    }
    atr[p]=t[j];
    for(p=k+1;p<n+2;p++,k++)
        atr[p]=t[k+1];
    for(j=0;j<n+1;j++)
    {
        if(atr[j]>atr[j+1])
            d[j]=atr[j]-atr[j+1];
        else
            d[j]=atr[j+1]-atr[j];
        sum+=d[j];
    }
    printf("\nAverage header movements:%f", (float)sum/n);
    getch();}
```

INPUT

Enter no.of tracks:9

Enter track position:55 58 60 70 18 90 150 160 184

OUTPUT

Tracks traversed Difference between tracks

150	50
160	10
184	24
90	94
70	20
60	10
58	2
55	3
18	37

Average header movements: 27.77

C) C-SCAN DISK SCHEDULING ALGORITHM

```
#include<stdio.h>
main()
{
    int t[20], d[20], h, i, j, n, temp, k, atr[20], tot, p, sum=0;
    clrscr();
    printf("enter the no of tracks to be traveresed");
    scanf("%d",&n);
    printf("enter the position of head");
    scanf("%d",&h);
    t[0]=0;t[1]=h;
    printf("enter      total      tracks");
    scanf("%d",&tot);
    t[2]=tot-1;
    printf("enter      the      tracks");
    for(i=3;i<=n+2;i++)
        scanf("%d",&t[i]);
    for(i=0;i<=n+2;i++)
        for(j=0;j<=(n+2)-i-1;j++)
            if(t[j]>t[j+1])
                {
                    temp=t[j];
                    t[j]=t[j+1];
                    t[j+1]=temp
                }
    for(i=0;i<=n+2;i++)
        if(t[i]==h);
        j=i;break;
        p=0;
        while(t[j]!=tot-1)
        {
            atr[p]=t[j];
            j++;
            p++;
        }
        atr[p]=t[j];
        p++;
        i=0;
        while(p!=(n+3) && t[i]!=t[h])
        {
            atr[p]=t[i]; i++;
            p++;
        }
}
```

```
for(j=0;j<n+2;j++)
{
    if(atr[j]>atr[j+1])
        d[j]=atr[j]-atr[j+1];
    else
        d[j]=atr[j+1]-atr[j];
    sum+=d[j];
}
printf("total header movements%d",sum);
printf("avg is %f", (float)sum/n);
getch();
}
```

INPUT

Enter the track position : 55 58 60 70 18 90 150 160 184
Enter starting position : 100

OUTPUT

Tracks traversed	Difference Between tracks
150	50
160	10
184	24
18	240
55	37
58	3
60	2
70	10
90	20

Average seek time : 35.7777779